



# Providing Differentiated Services, Congestion Management, and Deadlock Freedom in Dragonfly Networks

**Pedro Yébenes**, Jesús Escudero-Sahuquillo, Pedro J. García, Francisco J. Alfaro, Francisco J. Quiles

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## **Outline**

- Motivation
- Background
- Proposals description
- Evaluation
- Conclusions

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#### **Motivation**

#### **Interconnection Networks**

- Interconnection networks are key elements in HPC systems and datacenters.
  - Thousands of processing and/or storing nodes.
  - Applications need increasing computing power.
- The interconnection network may become the system **bottleneck** if not properly configured.

Achieving high network performance is mandatory.



Providing Differentiated Services, Congestion

Tianhe-2 (MilkyWay-2) 16000 nodes - Cores 3120000 TH-Express 2 **1st Top500** (November 2015)

March, 12th 2016

#### **Motivation**

#### **Interconnection Networks**

- Main design aspects of interconnection networks:
  - Topology
  - Routing Algorithm
  - Power consumption
  - Fault tolerance
  - Congestion control
  - Quality of service

#### **Motivation**

#### **Problem Statement**

- Minimal-path routing for Dragonfly networks is **not deadlock free** by default, requiring additional Virtual Channels (VCs) for deadlock freedom.
- Both congestion management and QoS can be provided by separating traffic flows into VCs.
- Thus, congestion management, QoS provision, and deadlock freedom require VCs for different purpose.
- There is not a joint and straightforward solution that offers these three functionalities at the same time.

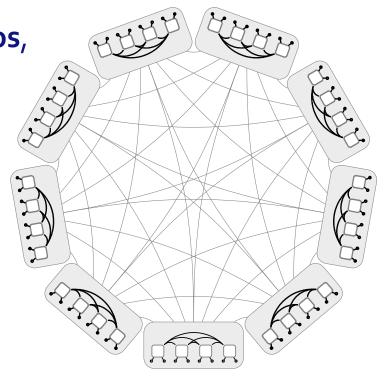
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## **Dragonfly Topology**

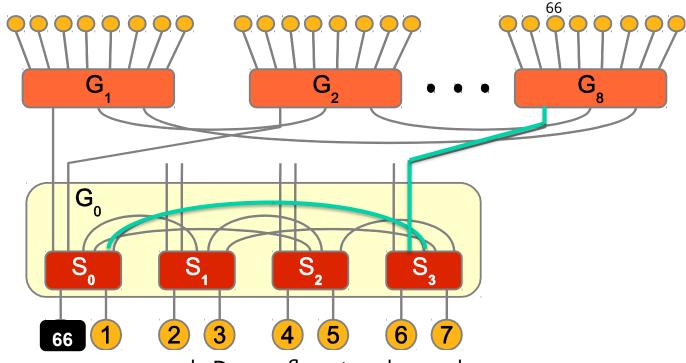
 Hierarchical high-performance topology consisting of a set of groups, each one composed of several switches where endnodes are attached.

 Low diameter, path diversity, high scalability, etc.



J. Kim, W. J. Dally, S. Scott, and D. Abts: **Technology-Driven, Highly-Scalable Dragonfly Topology**. SIGARCH 2008: vol. 36, pp. 77-88

## **Dragonfly Minimal-Path Routing**

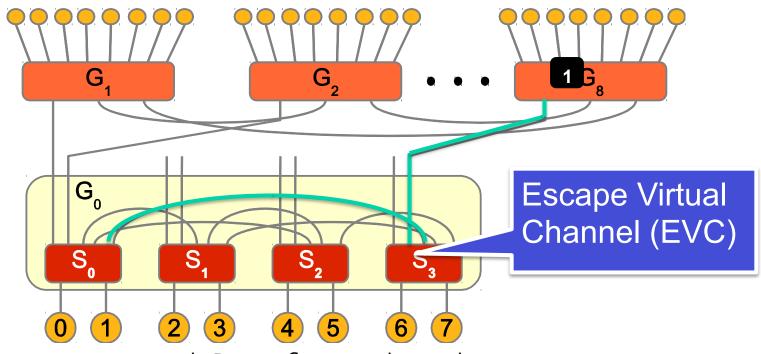


72-node Dragonfly network: a=4, h=2, p=2

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## **Dragonfly Minimal-Path Routing**



72-node Dragonfly network: a=4, h=2, p=2

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## Congestion

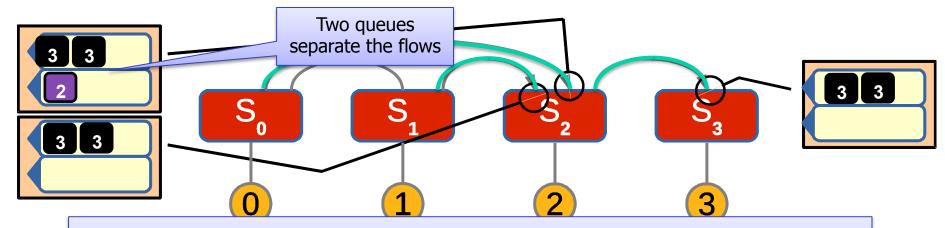
 Under congestion situations, network performance may degrade significantly.

 Head-of-Line blocking is the main problem derived from Congestion congestion. point **Congestion affects** also packets 3 belonging to other flows that are not contributing to congestion

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## **Queuing Schemes**

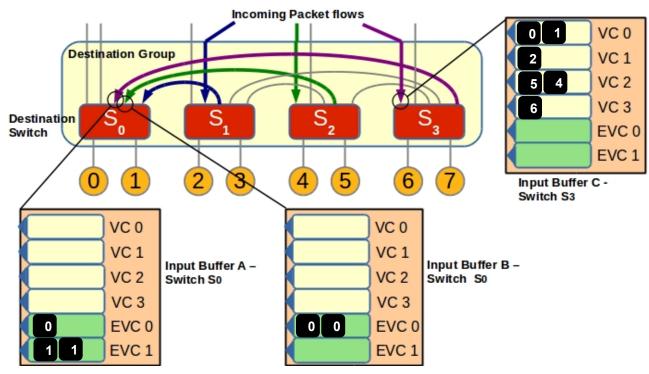
- Several queues, supporting Virtual Channels (VCs), or Virtual Lanes (VLs) are used at each port to separate traffic flows, reducing the HoL-blocking produced among them.
- A **static criterion** is used to map packets to queues.



The most efficient queuing schemes are tailored to a specific **network topology** and a specific **routing algorithm**.

#### Hierarchical 2-Level Queuing (H2LQ)

 Queuing scheme tailored to Dragonfly topology with MIN-path routing algorithm.



P. Yebenes, J. Escudero-Sahuquillo, P. J. Garcia and F. J. Quiles, "Efficient Queuing Schemes for HoL-Blocking Reduction in Dragonfly Topologies with Minimal-Path Routing," CLUSTER 2015 IEEE, pp. 817-824.

#### **Quality of Service**

- Usually, QoS is based on separating into different VCs
   traffic with different priorities or from different applications.
- Sometimes, VCs priorities are managed by using the Weighted Round Robin (WRR) algorithm, which is implemented by a weighted table.
- VCs with higher weight and/or more entries in the table have more priority.
   Weighted Table

VC	Weight
0	3
1	3
0	3
2	1

Total Weight VCo = 6/10 (60%) Total Weight VC1 = 3/10 (30%) Total Weight VC2 = 1/10 (10%)

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#### **Congestion Management + QoS**

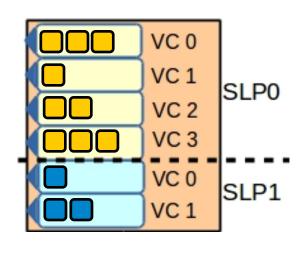
- CHADS: Combining HoL-blocking Avoidance and Differentiated Services.
- CHADS defines different **Service-Level Priorities** (SLPs) to identify the priority level of the applications.
- Each SLP is mapped to a **disjoint set of VCs**.
- A queuing scheme is used inside the set of VCs of the same SLP to prevent HoL blocking.
- Higher priority SLPs are mapped with more VCs.

P. Yebenes, J. Escudero-Sahuquillo, C. Gomez, P. J. Garcia F.J. Alfaro, and F. J. Quiles, "Combining Holblocking avoidance and differentiated services in high-speed interconnects," HiPC 2014

## **Congestion Management + QoS**

 CHADS: Combining HoL-blocking Avoidance and Differentiated Services.

SLP	VC	Weigth
	0	2
CLDo	1	2
SLPo	2	2
	3	2
SLP1	3	1
	4	1



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#### **Basic Ideas**

- Adapting CHADS to dragonfly networks.
- **SLPs** are also considered, each one assigned with different VCs.
- Congestion management inside each SLP by means of H2LQ.
- QoS provision by configuring Weighted Tables.
- **Deadlock freedom** using Escape VCs managed by different policies for configuring Escape Virtual Networks (EVNs):
  - Exclusive Escape Virtual Network (EEVN)
  - Common Exclusive Virtual Network (CEVN).

#### **Exclusive Escape Virtual Network (EEVN)**

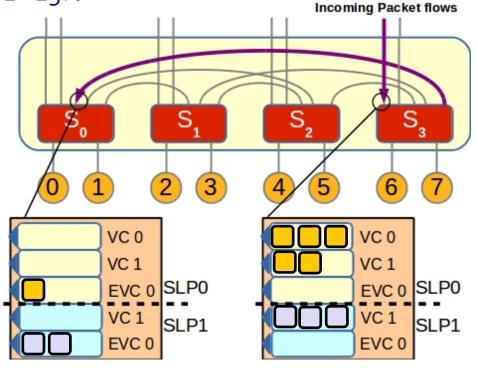
- Each SLP has a Standard Virtual Network (SVN) and an EVN.
- Packets use the SVN by default but are assigned to the EVN for avoiding deadlocks.
- Packets from different SLPs never interact.

#### **Exclusive Escape Virtual Network (EEVN)**

- 2 SLPs
  - Total Weight: SLP o = 75%, SLP1 = 25%
- 5 VCs

Weighted Table

SLP	VN	VC	Weigth
CLDo			2
SLPo	SVN	1	2
SLPo	EVN	2	2
SIP1	SVN	3	1
SLP1	EVN	4	1



#### **Common Escape Virtual Network (CEVN)**

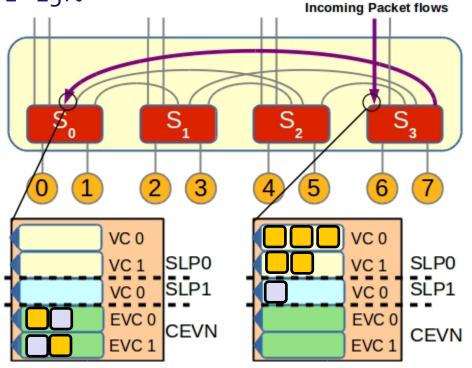
- Each SLP has a SVN.
- There is a single Common Escape VN (CEVN) shared by all the SLPs.
- Packets from different SLPs share VCs when they are in the CEVN.

#### **Common Escape Virtual Network (CEVN)**

- 2 SLPs
  - Total Weight: SLP o = 75%, SLP1 = 25%
- 5 VCs

Weighted Table

SLP	VN	VC	Weigth		
SLPo	SVN	0	3		
SLPO	2010	1	3		
SLP1	SVN	2	2		
SLP <sub>2</sub>	CEVN	3	1		
	CEVIN	4	1		



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#### **Simulation Tool**

#### OMNeT++-based simulator:

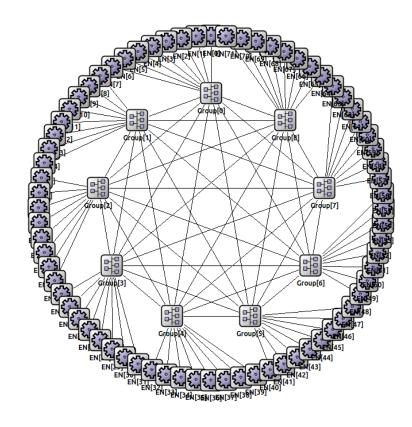
- Different topologies.
- Different routing algorithms.
- Different queuing schemes.
- Quality of Service support.



Pedro Yébenes, Jesús Escudero-Sahuquillo, Pedro J. García, Francisco J. Quiles: **Towards Modeling**Interconnection Networks of Exascale Systems with OMNeT++. PDP 2013

#### **Network Configurations**

• 4096-node dragonfly network (a=12, h=6, p=6).



#### **Traffic Patterns**

- 3 applications, each one assigned with a different SLP (SLPo, SLP1, SLP2), generating synthetic traffic at a rate of 70% of the link bandwidth with two traffic patterns:
  - Uniform traffic.
  - Zipf traffic:
    - Models traffic patterns with preferred destinations.
    - Traffic pattern similar to the ones produced by the collective communication schemes.

L. Breslau, Pei Cao, Li Fan, G. Phillips, and S. Shenker: **Web caching and Zipf-like distributions:** evidence and implications. INFOCOM '99: 126-134 vol.1

#### **EEVN VCs Configurations**

- Total weight per SLP in the Weighted Tables:
  - SLPo: 60%, SLP1: 30%, SL2: 10%
- Number of VCs for each SLP in EEVN (*EEVN-X* where *X* is the total number of required VCs).

SLPo		SLP1		SLP <sub>2</sub>		
Name	#SVCs	#EVCs	#SVCs	#EVCs	#SVCs	#EVCs
EEVN-54	12	6	12	6	12	6
EEVN-15	6	1	4	1	2	1
EEVN-8	3	1	1	1	1	1
EEVN-6	1	1	1	1	1	1

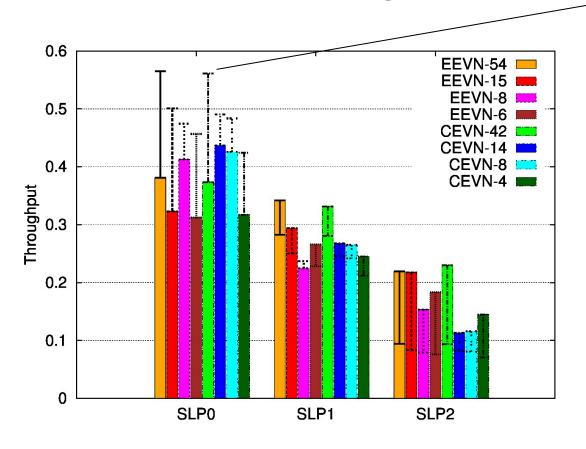
#### **CEVN VCs Configurations**

- Total weight per SLP in the Weighted Tables:
  - SLPo: 60%, SLP1: 30%, SL2: 10%
- Number of VCs for each SLP in CEVN (*CEVN-X* where *X* is the total number of required VCs).

Maron	SLPo	SLP1	SLP <sub>2</sub>	CEVN
Name	#SVC	#SVC	#SVC	#EVC
CEVN-42	12	12	12	6
CEVN-14	6	4	2	2
CEVN-8	3	2	1	2
CEVN-4	1	1	1	1

#### Results Uniform SLPo=70%, SLP1=70%, SLP2=70%

Metric: normalized throughput.

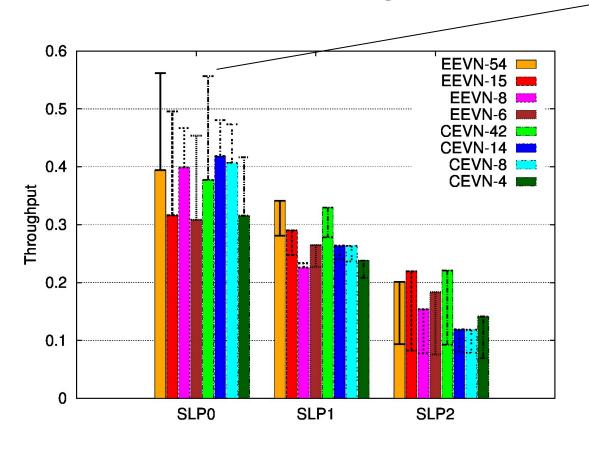


Distortion levels
show the expected
theoretical
throughput for a
given SLP,
according to the
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Providing Differentiated Services, Congestion

#### Results Zipf SLP0=70%, SLP1=70%, SLP2=70%

Metric: normalized throughput.



#### **Distortion levels** show the expected theoretical throughput for a given SLP, according to the network throughput and the WT configuration.

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#### **Conclusions**

#### **Advantages**

- CHADS technique has been updated for dragonfly networks.
- Differentiated services at network level, congestion management and deadlock freedom can be provided at the same time by means of EEVN and CEVN approaches.
- In general, CEVN is better than EEVN.
- The number of VCs configured has to be tightly with the weight configured in the Weighted Tables.

#### **Conclusions**

#### **Future directions**

- Analyzing these approaches with other routing algorithms suited to InfiniBand.
- Testing these approaches with other traffic patterns: application traces, adversarial, blocking collectives, etc.
- Exploring other configurations for Dragonfly networks.
- Exploring other approaches to better populate the weighted tables.





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